# A NEW CLASS OF *k*-RANK GRAPH ALGEBRAS, COMING FROM "*k*-CUBE GROUPS"

Sam A. Mutter with Aura-Cristiana Radu and Alina Vdovina

22nd February 2021

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This talk is based on:

S.A. Mutter, A.-C. Radu, and A. Vdovina

C\*-algebras of higher-rank graphs from groups acting on buildings, and explicit computation of their K-theory

https://arxiv.org/abs/2012.05561

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- 4. Higher-rank graph  $C^*$ -algebras
- 5. K-theory computations (non-scary version)

Why do we want more  $C^*$ -algebras?

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- Thus we have new bridges between operator algebras, geometric group theory, and category theory.

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- $\Lambda$  and d encode all the necessary information to replicate G (in fact we only need  $\Lambda^0$  and  $\Lambda^1$ ).



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#### DEFINITION (1-RANK GRAPH)

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Now, two morphisms  $\nu$ ,  $\mu$  are composable iff  $r(\nu) = s(\mu)$ ; then  $\mu \nu \in \Lambda$ .

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*We call the pair*  $(\Lambda, d)$  a 1-rank graph.

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- The converse construction also works.



Image: A matrix and a matrix

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- We say  $\Lambda$  is row-finite if  $\Lambda^{\mathbf{n}}(\nu)$  is finite, and that  $\Lambda$  has no sources if  $\Lambda^{\mathbf{n}}(\nu) \neq \emptyset$ , for all  $\nu \in \Lambda^{\mathbf{0}}$  and  $\mathbf{n} \in \mathbb{N}^k$ .

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Image: A matrix

EXAMPLES

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- So, in the 2-graph, we must define  $f_1e_1 = e_2f_1$ . Likewise, we identify  $f_2e_2 = e_1f_2$ .



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- Define range and domain maps  $r(\mathbf{m}, \mathbf{n}) := \mathbf{m}$ ,  $s(\mathbf{m}, \mathbf{n}) := \mathbf{n}$ , degree map  $d(\mathbf{m}, \mathbf{n}) := \mathbf{n} \mathbf{m}$ , and composition  $(\mathbf{l}, \mathbf{m})(\mathbf{m}, \mathbf{n}) := (\mathbf{l}, \mathbf{n})$ .

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#### EXAMPLES



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# *k*-cube groups

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• Let  $E_1, \ldots, E_k$  be alphabets, sets of even size  $\geq 4$ , each equipped with a fixed-point-free involution  $e \mapsto e^{-1}$ . E.g.  $E_1 = \{x_1, x_2, x_1^{-1}, x_2^{-1}\}$ .

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## k-CUBE GROUPS



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#### DEFINITION (k-CUBE GROUP)

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#### **DEFINITION** (*k*-CUBE GROUP)

Let  $E_1, \ldots, E_k$  be a collection of alphabets, and let R be the set of pointed squares labelled by (x, y, x', y') with  $x, x' \in E_i$  and  $y, y' \in E_j$  for i < j.

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- 2. The squares in R can be glued together to make pointed k-cubes.
- 3. Each combination  $\{e_1, \ldots, e_k \mid e_i \in E_i\}$  occurs at the corner of precisely one *k*-cube, up to symmetry.

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- 2. The squares in R can be glued together to make pointed k-cubes.
- 3. Each combination  $\{e_1, \ldots, e_k \mid e_i \in E_i\}$  occurs at the corner of precisely one *k*-cube, up to symmetry.

We call the group  $\Gamma := \langle E_1 \sqcup \cdots \sqcup E_k \mid xyx'y' = 1$  whenever  $(x, y, x', y') \in R \rangle$  a *k*-cube group.

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• What does it mean "up to symmetry"?

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• The symmetries map squares labelled from alphabets  $E_i$ ,  $E_j$  to other squares labelled from  $E_i$ ,  $E_j$ .

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#### EXAMPLES

# EXAMPLE (RUNGTANAPIROM, STIX, VDOVINA, 2019)

$$\begin{split} &\Gamma'_{\{3,5,7\}} := \langle a_1, a_2, b_1, b_2, b_3, c_1, c_2, c_3, c_4 \mid R \rangle, \ where \\ &R := \left\{ a_1 b_1 a_2 b_2, a_1 b_2 a_2 b_1^{-1}, a_1 b_3 a_2^{-1} b_1, a_1 b_3^{-1} a_1 b_2^{-1}, \\ &a_1 b_1^{-1} a_2^{-1} b_3, a_2 b_3 a_2 b_2^{-1}, a_1 c_1 a_2^{-1} c_2^{-1}, a_1 c_2 a_1^{-1} c_3, \\ &a_1 c_3 a_2^{-1} c_4^{-1}, a_1 c_4 a_1 c_1^{-1}, a_1 c_4^{-1} a_2 c_2, a_1 c_3^{-1} a_2 c_1, \\ &a_2 c_3 a_2 c_2^{-1}, a_2 c_4 a_2^{-1} c_1, c_1 b_1 c_3 b_3^{-1}, c_1 b_2 c_4 b_2^{-1}, c_1 b_3 c_4^{-1} b_2, \\ &c_1 b_3^{-1} c_4 b_3, c_1 b_2^{-1} c_2 b_1, c_1 b_1^{-1} c_4 b_1^{-1}, c_2 b_2 c_3^{-1} b_3^{-1}, \\ &c_2 b_3 c_4 b_1, c_2 b_3^{-1} c_3 b_3, c_2 b_2^{-1} c_3 b_2, c_2 b_1^{-1} c_3 b_1^{-1}, c_3 b_1 c_4 b_2 \right\} \end{split}$$

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k-CUBE GROUPS

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#### EXAMPLES

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EXAMPLES

#### EXAMPLE (k = 4, FREE GROUP)

Consider the product of four free groups, each with two generators, defined as follows:

$$\mathbb{F}_2^4 := \langle a_1, a_2, b_1, b_2, c_1, c_2, d_1, d_2 \mid [a_i, b_j], [a_i, c_j], [a_i, d_j], \ [b_i, c_j], [b_i, d_j], [c_i, d_j] \textit{ for all } i, j \in \{1, 2\} 
angle.$$

This is a 4-cube group.

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EXAMPLES

#### EXAMPLE (k = 4, MRV)

 $\Gamma_{\{1,2,3,4\}} := \langle a_1, a_2, a_3, b_1, b_2, b_3, c_1, c_2, c_3, d_1, d_2, d_3 \mid R \rangle$ , where  $R := \{a_1b_1a_3^{-1}b_1, a_1b_1^{-1}a_2^{-1}b_3, a_1b_2a_2b_2, a_1b_2^{-1}a_3b_3^{-1}, a_1b_3a_2^{-1}b_1^{-1}, a_1b_3^{-1}a_3b_2^{-1}, a_1b_3a_2^{-1}b_1^{-1}, a_1b_3^{-1}a_3b_2^{-1}, a_1b_3a_2^{-1}b_1^{-1}, a_1b_3a_2^{-1}b_1^{-1}a_3b_2^{-1}, a_1b_3a_2^{-1}b_1^{-1}a_3b_2^{-1}, a_1b_3a_2^{-1}b_1^{-1}a_3b_2^{-1}a_3b_2^{-1}, a_1b_3a_2^{-1}b_1^{-1}a_3b_2^{-1}a_3b_2^{-1}a_3b_3^{-1}a_3b_3^{-1}a_3b_3^{-1}a_3b_2^{-1}a_3b_3^{-1}a_3b_$  $a_{2}b_{3}a_{3}b_{3}, a_{3}b_{1}a_{2}^{-1}b_{2}, a_{3}b_{2}a_{2}^{-1}b_{1}, a_{1}c_{1}a_{2}^{-1}c_{1}, a_{1}c_{1}^{-1}a_{1}c_{2}^{-1}, a_{1}c_{2}a_{1}^{-1}c_{2}^{-1}, a_{1}c_{3}a_{3}c_{3}, a_{1}c_{2}a_{1}^{-1}c_{2}^{-1}, a_{1}c_{3}a_{3}c_{3}, a_{1}c_{2}a_{1}^{-1}c_{2}^{-1}, a_{1}c_{3}a_{3}c_{3}, a_{1}c_{3}c_{3}a_{3}c_{3}, a_{1}c_{3}c_{3}c_{3}, a_{1$  $a_2c_1a_2c_2^{-1}, a_2c_2a_3^{-1}c_2, a_2c_3a_2^{-1}c_3^{-1}, a_3c_1^{-1}a_3^{-1}c_1, a_3c_2a_3c_3^{-1}, a_1d_1a_3^{-1}d_3, a_1d_1^{-1}a_2d_2,$  $a_1d_2a_2d_1^{-1}, a_1d_2^{-1}a_1d_2^{-1}, a_1d_3a_2^{-1}d_1, a_2d_1a_2d_2^{-1}, a_2d_2^{-1}a_3d_3, a_2d_3a_3d_2^{-1}, a_3d_1a_3d_2,$  $b_1c_1b_3^{-1}c_1, b_1c_1^{-1}b_2^{-1}c_3, b_1c_2b_2c_2, b_1c_2^{-1}b_3c_2^{-1}$  $b_1c_3b_2^{-1}c_1^{-1}, b_1c_3^{-1}b_3c_2^{-1}, b_2c_3b_3c_3, b_3c_1b_2^{-1}c_2, b_3c_2b_2^{-1}c_1,$  $\begin{array}{c} b_1d_1b_2^{-1}d_1, b_1d_1^{-1}b_1d_3^{-1}, b_1d_2b_1^{-1}d_2^{-1}, b_1d_3b_3d_3, \\ b_2d_1b_2d_2^{-1}, b_2d_2b_3^{-1}d_2, b_2d_3b_2^{-1}d_3^{-1}, b_3d_1^{-1}b_3^{-1}d_1, b_3d_2b_3d_3^{-1}, \end{array}$  $c_1d_1c_2^{-1}d_1, c_1d_1^{-1}c_2^{-1}d_3, c_1d_2c_2d_2, c_1d_2^{-1}c_3d_2^{-1},$  $c_1d_3c_2^{-1}d_1^{-1}, c_1d_3^{-1}c_3d_2^{-1}, c_2d_3c_3d_3, c_3d_1c_2^{-1}d_2, c_3d_2c_2^{-1}d_1$ 

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#### EXAMPLES



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Given a k-cube group  $\Gamma = \langle E_1, \ldots, E_k | \mathcal{R} \rangle$ , the subgroup generated by alphabets  $E_1, \ldots, E_{k-1}$  is a (k-1)-cube group.

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We can find k different (k - 1)-cube subgroups in this way. The group  $\Gamma$  is the product of these groups, amalgamated over the free groups generated by their pairwise intersections.

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• In general the converse is not true—it is difficult to find a family of k-cube groups whose amalgamated product forms a (k + 1)-cube group.

**CONSTRUCTING A SHIFT SYSTEM** 

• Let  $\Delta$  be the rank k affine building which is the k-dimensional cube complex  $T_{|E_1|} \times \cdots \times T_{|E_k|}$ .

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**CONSTRUCTING A SHIFT SYSTEM** 

Let Δ be the rank k affine building which is the k-dimensional cube complex T<sub>|E<sub>1</sub>|</sub> × · · · × T<sub>|E<sub>k</sub>|</sub>. We identify elements of Γ with edges of Δ, so the set of k-cubes S<sub>k</sub> can be identified with the set of pointed, oriented chambers of Δ.

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**DEFINITION (ADJACENCY FUNCTIONS)** 

#### CONSTRUCTING A SHIFT SYSTEM

• Let  $\Delta$  be the rank k affine building which is the k-dimensional cube complex  $T_{|E_1|} \times \cdots \times T_{|E_k|}$ . We identify elements of  $\Gamma$  with edges of  $\Delta$ , so the set of k-cubes  $S_k$  can be identified with the set of pointed, oriented chambers of  $\Delta$ .

### **DEFINITION (ADJACENCY FUNCTIONS)**

Let  $A, B \in S_k$ . We define adjacency functions  $M_1, \ldots, M_k : S_k \times S_k \rightarrow \{0, 1\}$ , where  $M_i(A, B) = 1$  iff:

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• The first (k - 1)-dimensional face of A labelled by alphabets  $E_1, \ldots, \hat{E}_i, \ldots, E_k$  coincide with the second such (k - 1)-face of B, and

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- The first (k 1)-dimensional face of A labelled by alphabets  $E_1, \ldots, \hat{E}_i, \ldots, E_k$  coincide with the second such (k 1)-face of B, and
- Whenever  $M_i(A, B) = 1$ , if we stack A and B together so that their common (k-1)-faces overlap, we never have e and  $e^{-1}$  pointing to the same vertex.

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### **CONSTRUCTING A SHIFT SYSTEM**

• Consider this example for k = 3.



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#### CONSTRUCTING A SHIFT SYSTEM

- Consider this example for k = 3.
- Suppose all black edges are labelled from *E*<sub>1</sub>, magenta from *E*<sub>2</sub>, and blue from *E*<sub>3</sub>.
- $M_2(A,B) = 1$  iff  $(u_1, v_1, u_2, v_2) = (a_4^{-1}, b_3^{-1}, a_3^{-1}, b_2^{-1})$  and  $c_1 \neq w_1^{-1}$ ,  $c_3 \neq w_3^{-1}$ , etc.



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- $M_3(A,C) = 1$  iff  $(u_2, w_2, u_3, w_3) = (x_1^{-1}, z_4^{-1}, x_4^{-1}, z_1^{-1})$  and  $y_3 \neq v_3^{-1}$ , etc.



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A k-cube group induces a k-rank graph

## THEOREM (MRV)

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A k-cube group induces a k-rank graph

### THEOREM (MRV)

Let  $\Gamma$  be a k-cube group, and write the adjacency functions  $M_1, \ldots, M_k$  as matrices. Then each  $M_i$  has at least 3 non-zero entries in each row, and  $M_iM_j = M_jM_i$  for all i, j.

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### THEOREM (MRV)

A k-cube group  $\Gamma$  has the unique common extension property.

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### A k-cube group induces a k-rank graph

 The UCE property says that in every dimension 2 ≤ n ≤ k, if we start with an n-dimensional cube and find n adjacent n-cubes (one in each direction), then we can uniquely extend the arrangement into a 2 × 2 supercube.



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- Assign to each k-cube in  $\Gamma$  a vertex.
- Draw an *i*-coloured arrow from A to B if B is adjacent to A in the *i* direction.
- The UCE property ensures that this *k*-coloured graph has all the factorisation properties of a *k*-rank graph.

![](_page_123_Figure_6.jpeg)

The k-rank graph algebra

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SAM A. MUTTER

The k-rank graph algebra

### DEFINITION (k-RANK GRAPH ALGEBRA OF KUMJIAN-PASK)

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The k-rank graph algebra

DEFINITION (k-RANK GRAPH ALGEBRA OF KUMJIAN-PASK)

Let  $\Lambda = (\Lambda, d)$  be a row-finite k-rank graph with no sources.

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The k-rank graph algebra

## DEFINITION (*k*-RANK GRAPH ALGEBRA OF K<u>UMJIAN-PASK</u>)

Let  $\Lambda = (\Lambda, d)$  be a row-finite k-rank graph with no sources. We define the k-rank graph algebra  $\mathcal{A}(\Lambda)$  to be the universal C<sup>\*</sup>-algebra generated by a family  $\{s_{\lambda} \mid \lambda \in \Lambda\}$  of partial isometries which have the following properties:

The k-rank graph algebra

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1. The set  $\{s_{\nu} \mid \nu \in \Lambda^{\mathbf{0}}\}$  satisfies  $(s_{\nu})^2 = s_{\nu} = s_{\nu}^*$  and  $s_u s_{\nu} = 0$  for all  $u \neq \nu$ .

2. If 
$$r(\lambda) = s(\mu)$$
 for some  $\lambda, \mu \in \Lambda$ , then  $s_{\mu\lambda} = s_{\mu}s_{\lambda}$ .

3. For all  $\lambda \in \Lambda$ , we have  $s_{\lambda}^* s_{\lambda} = s_{s(\lambda)}$ .

4. For all vertices  $v \in \Lambda^{\mathbf{0}}$  and  $\mathbf{n} \in \mathbb{N}^k$ , we have:  $s_v = \sum_{\lambda \in \Lambda^{\mathbf{n}}(v)} s_\lambda s_\lambda^*$ .

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The k-rank graph algebra

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- 4. For all vertices  $v \in \Lambda^{\mathbf{0}}$  and  $\mathbf{n} \in \mathbb{N}^k$ , we have:  $s_v = \sum_{\lambda \in \Lambda^{\mathbf{n}}(v)} s_\lambda s_\lambda^*$ .

Without the row-finiteness condition, property 4 is not well-defined.

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THEOREM (EVANS, 2008)

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### THEOREM (EVANS, 2008)

Let  $\Lambda$  be a row-finite 2-rank graph with no sources. Then the K-theory of  $\mathcal{A}(\Lambda)$  is well understood.

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Let 
$$k = 3$$
, and  $\partial_1 := [I - M_1^T, I - M_2^T, I - M_3^T]$ ,  $\partial_3 := [I - M_3^T, M_2^T - I, I - M_1^T]^T$ ,

$$egin{aligned} \partial_2 &:= egin{bmatrix} M_2^T - I & M_3^T - I & 0 \ I - M_1^T & 0 & M_3^T - I \ 0 & I - M_1^T & I - M_2^T \end{bmatrix} \end{aligned}$$

Then  $K_1(\mathcal{A}(\Lambda)) \cong \ker(\partial_1) / \operatorname{im}(\partial_2) \oplus G_1$ , and there is a short exact sequence

$$0 \longrightarrow \operatorname{coker}(\partial_1)/G_0 \longrightarrow K_0(\mathcal{A}(\Lambda)) \longrightarrow \operatorname{ker}(\partial_2)/\operatorname{im}(\partial_3) \longrightarrow 0,$$

where  $G_0 \subseteq \operatorname{coker}(\partial_1)$  and  $G_1 \subseteq \ker(\partial_3)$ .

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### THEOREM (MRV)

There are similar (nastier) short exact sequences and isomorphisms for k = 4 and k = 5.

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The k-rank graph algebra

### THEOREM (MRV)

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### THEOREM (MRV)

Let  $\Gamma$  be a k-cube group, and  $\Lambda$  be its induced k-rank graph. Then  $\mathcal{A}(\Lambda)$  is a *Kirchberg algebra*— in particular, it is completely classified (up to isomorphism) by its K-theory.

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**COMPUTATION OF K-THEORY** 

## EXAMPLE (RUNGTANAPIROM, STIX, VDOVINA, 2019)

$$\Gamma'_{\{3,5,7\}} := \langle a_1, a_2, b_1, b_2, b_3, c_1, c_2, c_3, c_4 \mid R \rangle.$$

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## EXAMPLE (RUNGTANAPIROM, STIX, VDOVINA, 2019)

$$\begin{split} &\Gamma'_{\{3,5,7\}} := \langle a_1, a_2, b_1, b_2, b_3, c_1, c_2, c_3, c_4 \mid R \rangle. \\ & We \ have \ K_1(\mathcal{A}(\Gamma)) \cong \mathbb{Z}^{21} \oplus (\mathbb{Z}/2)^6 \oplus (\mathbb{Z}/4)^2 \oplus (\mathbb{Z}/12)^2 \oplus G_1 \ and \end{split}$$

$$0 \longrightarrow \frac{\mathbb{Z}^7 \oplus (\mathbb{Z}/4)^2 \oplus (\mathbb{Z}/12)}{G_0} \longrightarrow K_0(\mathcal{A}(\Gamma)) \longrightarrow \mathbb{Z}^{21} \oplus (\mathbb{Z}/4)^2 \oplus (\mathbb{Z}/12) \longrightarrow 0,$$

for some  $G_0 \subseteq \mathbb{Z}^7 \oplus (\mathbb{Z}/4)^2 \oplus (\mathbb{Z}/12)$  and  $G_1 \subseteq \mathbb{Z}^7$ .

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for some  $G_0 \subseteq \mathbb{Z}^7 \oplus (\mathbb{Z}/4)^2 \oplus (\mathbb{Z}/12)$  and  $G_1 \subseteq \mathbb{Z}^7$ .

We can deduce that the torsion-free part of  $K_0$  is isomorphic to  $\mathbb{Z}^r$ , and that  $K_1 \cong \mathbb{Z}^r \oplus (\mathbb{Z}/2)^6 \oplus (\mathbb{Z}/4)^2 \oplus (\mathbb{Z}/12)^2$ , for some  $21 \le r \le 28$ .

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## EXAMPLE (k = 4, FREE GROUP)

$$\begin{split} \mathbb{F}_2^4 &:= \langle a_1, a_2, b_1, b_2, c_1, c_2, d_1, d_2 \mid [a_i, b_j], [a_i, c_j], [a_i, d_j], \\ [b_i, c_j], [b_i, d_j], [c_i, d_j] \textit{ for all } i, j \in \{1, 2\} \rangle. \end{split}$$

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Our short exact sequences give us (something like)  $K_0(\mathcal{A}(\mathbb{F}_2^4)) \cong K_1(\mathcal{A}(\mathbb{F}_2^4)) \cong \mathbb{Z}^{64} \oplus \mathbb{Z}^r$ , where  $0 \le r \le 64$ .

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Compare our sequences with the values obtained for  $K_0$  and  $K_1$  by the Künneth Theorem for tensor products:  $K_0 \cong K_1 \cong \mathbb{Z}^{128}$  and we see that r is maximal.

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**COMPUTATION OF K-THEORY** 

### EXAMPLE (MRV)

 $\Gamma_{\{1,2,3,4\}} := \langle a_1, a_2, a_3, b_1, b_2, b_3, c_1, c_2, c_3, d_1, d_2, d_3 \mid R \rangle.$ 

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## EXAMPLE (MRV)

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We have written a program in Python which determines whether a group is a 4-cube group, and if so, outputs four adjacency matrices. In this example,  $\Gamma$  is a 4-cube group, but the adjacency matrices are very large. We are still awaiting the computations of the K-theory in MAGMA.

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## **Epilogue**

- The K-theory of these algebras is a great way to distinguish them from each other, showing that our constructions do indeed produce new C\*-algebras.
- Computing K-theory of k-rank graph algebras is hard (e.g. SAM 2020).
- Next step: make the code work!

S.A. Mutter, A.-C. Radu, and A. Vdovina, C\*-algebras of higher-rank graphs from groups acting on buildings, and explicit computation of their K-theory https://arxiv.org/abs/2012.05561

S.A. Mutter, The K-theory of the  $C^*$ -algebras of 2-rank graphs associated to complete bipartite graphs

https://arxiv.org/abs/2004.11602

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